Titanium Performance and Potential: an NPB Experimental Study

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Take-Home Messages

Titanium:

- allows for elegant and concise programs
- gets comparable performance to Fortran+MPI on three common yet diverse scientific kernels (NPB)
- is well-suited to real-world applications
- is portable (runs everywhere)



NAS Parallel Benchmarks

Conjugate Gradient (CG)

- *Computation*: Mostly sparse matrix-vector multiply (SpMV)
- Communication: Mostly vector and scalar reductions

• 3D Fourier Transform (FT)

- *Computation*: 1D FFTs (using FFTW 2.1.5)
- Communication: All-to-all transpose
- Multigrid (MG)
 - *Computation*: 3D stencil calculations
 - Communication: Ghost cell updates



Titanium Overview

- Titanium is a Java dialect for parallel scientific computing
 - No JVM, no JIT, and no dynamic class loading

Titanium is extremely portable

- Ti compiler is source-to-source, and first compiles to C for portability
- Ti programs run everywhere- uniprocessors, shared memory, and distributed memory systems

All communication is one-sided for performance

GASNet communication system (not MPI)



Presented Titanium Features

Features in addition to standard Java:

- Flexible and efficient multi-dimensional arrays
- Built-in support for multi-dimensional domain calculus
- Partitioned Global Address Space (PGAS) memory model
- Locality and sharing reference qualifiers
- Explicitly unordered loop iteration
- User-defined immutable classes
- Operator-overloading
- Efficient cross-language support
- Many others not covered...



Titanium Arrays

• Ti Arrays are created and indexed using *points*:

 double [3d] gridA = new double [[-1,-1,-1]:[256,256,256]];

 (MG)
 Lower Bound

- gridA has a rectangular index set (*RectDomain*) of all points in box with corners [-1, -1, -1] and [256, 256, 256]
- Points and RectDomains are first-class types
- The power of Titanium arrays lies in:
 - · Generality: indices can start at any point
 - · Views: one array can be a subarray of another



Foreach Loops

 Foreach loops allow for unordered iterations through a RectDomain:

```
public void square(double [3d] gridA, double [3d] gridB) {
   foreach (p in gridA.domain()) {
     gridB[p] = gridA[p] * gridA[p];
   }
}
```

- These loops:
 - allow the compiler to reorder execution to maximize performance
 - require only one loop even for multidimensional arrays
 - avoid off-by-one errors common in *for* loops





Point Operations

• Titanium allows for arithmetic operations on Points:

This makes the MG stencil code more readable and concise





Titanium Parallelism Model

Ti uses an SPMD model of parallelism

- Number of threads is fixed at program startup
- Barriers, broadcast, reductions, etc. are supported
- Programmability using a Partitioned Global Address Space (i.e., direct reads and writes)
 - Programs are portable across shared/distributed memory
 - Compiler/runtime generates communication as needed
 - User controls data layout locality; key to performance





PGAS Memory Model

- Global address space is logically partitioned
 - Independent of underlying hardware (shared/distributed)
 - Data structures can be spread over partitions of shared space
- References (pointers) are either local or global (meaning possibly remote)





Distributed Arrays

 Titanium allows construction of distributed arrays in the shared Global Address Space:



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Domain Calculus and Array Copy

- Full power of Titanium arrays combined with PGAS model
- Titanium allows set operations on RectDomains:

// update overlapping ghost cells of neighboring block
data[neighborPos].copy(myData.shrink(1));
(MG)

- The copy is only done on intersection of array RectDomains
- Titanium also supports nonblocking array copy





The Local Keyword and Compiler Optimizations

 Local keyword ensures that compiler statically knows that data is local:

double [3d] myData = (double [3d] local) data[myBlockPos];

- This allows the compiler to use more efficient native pointers to reference the array
 - Avoid runtime check for local/remote
 - Use more compact pointer representation
- Titanium optimizer can often automatically propagate locality info using *Local Qualifier Inference* (LQI)



Is LQI (Local Qualifier Inference) Useful?



- LQI does a solid job of propagating locality information
- Speedups:
 - CG- 58%
 improvement
 - MG- 77% improvement



Immutable Classes

For small objects, would sometimes prefer:

- to avoid level of indirection and allocation overhead
- to pass by value (copying of entire object)
- especially when immutable (fields never modified)
 - · Extends idea of primitives to user-defined data types
- Example: Complex number class





Operator Overloading

For convenience, Titanium allows operator overloading

- Overloading in Complex makes the FT benchmark more readable
- Similar to operator overloading in C++

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Cross-Language Calls

- Titanium supports efficient calls to kernels/libraries in other languages
 - no data copying required
- Example: the FT benchmark calls the FFTW library to perform the local 1D FFTs
- This encourages:
 - shorter, cleaner, and more modular code
 - the use of tested, highly-tuned libraries



Are these features expressive?



- Compared line counts of timed, uncommented portion of each program
 - MG and FT disparities mostly due to Ti domain calculus and array copy
- CG line counts are similar since Fortran version is already compact



Testing Platforms

Opteron/InfiniBand (NERSC / Jacquard):

- *Processor*: Dual 2.2 GHz Opteron (320 nodes, 4 GB/node)
- *Network*: Mellanox Cougar InfiniBand 4x HCA

• G5/InfiniBand (Virginia Tech / System X):

- Processor: Dual 2.3 GHz G5 (1100 nodes, 4 GB/node)
- *Network*: Mellanox Cougar InfiniBand 4x HCA



Problem Classes

	Matrix or Grid Dimensions	Iterations
CG Class C	150,000 ²	75
CG Class D	1,500,000 ²	100
FT Class C	512 ³	20
MG Class C	512 ³	20
MG Class D	1024 ³	50

All problem sizes shown are relatively large



Data Collection and Reporting

- Each data point was run three times, and the minimum of the three is reported
- For a given number of procs, the Fortran and Titanium codes were run on the same nodes (for fairness)
- All the following speedup graphs use the *best* time at the lowest number of processors as the baseline for the speedup





FT Speedup





- All versions of the code use FFTW 2.1.5 for the serial 1D FFTs
- Nonblocking array copy allows for comp/comm overlap
- Max Mflops/proc:

	For	Ti(bl)	Ti(nbl)
G5	238	294	350
Opt.	203	268	318





MG Speedup



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CG Speedup



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Other Applications in Titanium

Larger Applications

- Heart and cochlea simulations (E. Givelberg, K. Yelick, A. Solar-Lezama, J. Su)
- AMR Elliptic PDE solver (P. Colella, T. Wen)

Other Benchmarks and Kernels

- Scalable Poisson solver for infinite domains
- Unstructured mesh kernel: EM3D
- Dense linear algebra: LU, MatMul
- Tree-structured n-body code
- Finite element benchmark



Conclusions

Titanium:

- Captures many abstractions needed for common scientific kernels
- Allows for more productivity due to fewer lines of code
- Performs comparably and sometimes better to Fortran w/MPI
- Provides more general distributed data layouts and irregular parallelism patterns for real-world problems (e.g., heart simulation, AMR)

